## **Chapter 2 Homework**

- 1. [10 points] Do Problem 2.3(a) in the text (page 71), but start with the recurrence T(n) = 3T(n/2) + n.
- (a) [4 points] What is the general kth term in this case?
- (b) [3 points] What value of k should be plugged in to get the answer?
- (c) [3 points] What is the resulting solution of the recurrence? Do not use the Master Theorem, but use the answers to parts (a) and (b) of this question.
- **2.** [20 points] Solve each of the following recurrences, giving a  $\Theta$ -bound.
- (a)  $T(n) = 3T(n/3) + n^2$
- (b)  $T(n) = 27T(n/3) + n^2$
- (c)  $T(n) = 9T(n/3) + n^2$
- (d)  $T(n) = T(n-1) + n^2$ .

**Hint:** The Master Theorem doesn't apply here. Do a repeated substitution (as in Problem 1 above), maybe with a small value of n (say, n = 5) to see what's going on.

- **3.** [10 points] You are given an infinite array  $a[\cdot]$ . You are *not* given the value of n. Describe an algorithm FIND(x, a), where
  - x is an integer,
  - a is an infinite array, in which the first n cells contain integers in sorted order and the other cells are marked with  $\infty$ ,

and whose return value is

$$FIND(x, a) = \begin{cases} j & \text{if } x = a[j] \text{ for some index } j, \\ \infty & \text{otherwise.} \end{cases}$$

The running time of this algorithm must be  $O(\log n)$ .

**Hint:** Note that you do not know the value of n. So, FIND's first task is determining a good upper bound on the value of n.